

Comments on *Judgment Aggregation without Paradox*

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 - **Connection**: applying majority-voting on isolated premises blinds the procedure to the logical connections between them. [Discussed in section 2]

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 - **Connection**: applying majority-voting on isolated premises blinds the procedure to the logical connections between them. [Discussed in section 2]
- One more common drawback of PBP & Fusion
Manipulability: aggregated outcome can be manipulated by strategic voting. [Discussed in section 3]

The San Francisco Art School is looking for models (of any sex) for their painting classes. Jerry shows up for the auditions. A seven-member committee votes on the following:

- A: Jerry is attractive
- B: Jerry is poor in social skills
- C: Jerry is accepted
- **Integrity Constraint:** $((A \& B) \vee (\sim A \& \sim B)) \equiv C$

	A	B	C	$A \& B$	$\sim A \& \sim B$	$A \equiv B$
Voter 1	1	1	1	1	0	1
Voter 2	1	1	1	1	0	1
Voter 3	1	1	1	1	0	1
Voter 4	1	0	0	0	0	0
Voter 5	1	0	0	0	0	0
Voter 6	0	1	0	0	0	0
Voter 7	0	1	0	0	0	0
Majority	1	1	0	0	0	0

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- Example of *instability* of PBP: if we identify A and B as premises, PBP's outcome on C is 1.
- If we collect majority on the complex propositions $A \& B$ and $\sim A \& \sim B$ as premises, PBP's outcome on C is 0.

Applying the Fusion Procedure

- $Mod(\mathcal{K}_1) = Mod(\mathcal{K}_2) = Mod(\mathcal{K}_3) = \{(1, 1, 1)\}$
- $Mod(\mathcal{K}_4) = Mod(\mathcal{K}_5) = \{(1, 0, 0)\}$
- $Mod(\mathcal{K}_6) = Mod(\mathcal{K}_7) = \{(0, 1, 0)\}$
- $Mod(IC) = \{(1, 1, 1), (0, 0, 1), (1, 0, 0), (0, 1, 0)\}$

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A	B	C	\mathcal{K}_1	\mathcal{K}_2	\mathcal{K}_3	\mathcal{K}_4	\mathcal{K}_5	\mathcal{K}_6	\mathcal{K}_7	$d_{\Sigma}(I, E)$
1	1	1	0	0	0	2	2	2	2	8
1	0	0	2	2	2	0	0	2	2	10
0	1	0	2	2	2	2	2	0	0	10
0	0	1	2	2	2	2	2	2	2	14

A Language Variation

- Let us now move to a language that does not contain A , and B but where there are atomic sentences $D :\approx (A \& B)$ and $E :\approx (\sim A \& \sim B)$.

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- Let us now move to a language that does not contain A , and B but where there are atomic sentences $D \approx (A \& B)$ and $E \approx (\sim A \& \sim B)$.
- Same vote as before, but now the integrity constraints are expressed in terms of D , E , and C .
- **IC:** $(D \vee E) \equiv C$ and $D \rightarrow \sim E$.
 - Note the extra-IC: it takes care of the connection between $(A \& B)$ and $(\sim A \& \sim B)$

Reversal!

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- Recall: $D : \approx (A \& B)$ and $E : \approx (\sim A \& \sim B)$.
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D	E	C	\mathcal{K}_1	\mathcal{K}_2	\mathcal{K}_3	\mathcal{K}_4	\mathcal{K}_5	\mathcal{K}_6	\mathcal{K}_7	$d_{\Sigma}(I, E)$
1	0	1	0	0	0	2	2	2	2	8
0	1	1	2	2	2	2	2	2	2	14
0	0	0	2	2	2	0	0	0	0	6

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- Point of difference with instability of PBP:
 - (Proper) instability of PBP: different ways of applying the same procedure yield different outcomes—holding fixed the identification of the atoms.
 - In the current case, we need to carry out the fusion procedure exactly in the same way, but look at two different languages. (*linguistic instability*).
- The common feature between linguistic instability and instability proper is that in both cases the aggregation procedure sees a difference where there should be none.

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- Let \mathcal{L}_2 have three sentences X, Y, Z , such that: $Z : \approx C$,
 $Y : \approx B \equiv C$, $X : \approx A \equiv B$.

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 - have the same number of atomic sentences
 - in virtue of the expressive equivalence, there is never a need for extra-integrity constraints.
- Let \mathcal{L}_2 have three sentences X, Y, Z , such that: $Z \approx C$, $Y \approx B \equiv C$, $X \approx A \equiv B$.
- E.g.: A vote on $(A \& B) \equiv C$ becomes a vote on $[(X \equiv (Z \equiv Y)) \& (Z \equiv Y)] \equiv Z$.

My conclusion in this section is that, to defend her approach, and at the same time be able to use *instability* as an objection to PBP, Gabriella owes both of the following:

- 1 An explanation of why *linguistic* instability is harmless for the Fusion procedure.

My conclusion in this section is that, to defend her approach, and at the same time be able to use *instability* as an objection to PBP, Gabriella owes both of the following:

- 1 An explanation of why *linguistic* instability is harmless for the Fusion procedure.
- 2 An account of why the explanation in (1) cannot be used by a supporter of PBP against the instability criticism.

- Distinction: (i) collective judgments in which the outcome must be supported by collectively accepted reasons vs. (ii) judgments in which it is enough to reach some collective conclusion.

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- Standard view:
 - (i) is best dealt by procedures which, like PBP and fusion, give a complete outcome
 - (ii) is better dealt by the Conclusion Based Procedure (CPB).
- In the art-school case, we can imagine contexts in which the reasons for acceptance of C are important (e.g. if Jerry is accepted because attractive and unapproachable ($A \& B$) then he gets a worse contract than he would get if he were accepted because unattractive and easygoing ($\sim A \& \sim B$).

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- But here, the fusion procedure recommends the same outcome as PBP, even though only a minority supports $(A \& B)$ (and thus $(A \equiv B)$).
- **Question:** in what sense does the Fusion procedure do better when it comes to keeping track of logical connections between premises?

Independence on a Set of Propositions

- F. Dietrich and C. List, *Strategy-proof judgment aggregation*, forthcoming.

(*Ind_Y*) An aggregation function F is independent (on a set of propositions Y) iff for every proposition $p \in Y$, and $n \in \omega$ there is a function $\phi : \{0, 1\}^n \rightarrow \{0, 1\}$ such that $p \in F(\mathcal{K}_1, \dots, \mathcal{K}_m)$ iff $\phi(v_1(p), \dots, v_n(p)) = 1$ where $v_i(p)$ is \mathcal{K}_i 's vote on p .

- For each member p of Y the aggregated outcome on p is a function of the pattern of individual judgments on p .

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- For each member p of Y the aggregated outcome on p is a function of the pattern of individual judgments on p .
- In our example, whenever $Y \supset \{C\}$, both *PBP* and *Fusion* are not Independent on Y .

	A	B	C		A	B	C
Voter 1	1	1	1		1	1	1
Voter 2	1	1	1		1	1	1
Voter 3	1	1	1		1	1	1
Voter 4	1	0	0		1	0	0
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Voter 6	0	1	0		1	0	0
Voter 7	0	1	0		1	0	0
PBP + Fusion	1	1	1		1	0	0

- A cost of relaxing independence: some voters can manipulate the outcome on some members of Y (in this case on the conclusion C) by strategic voting.

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- More formally, let F be an aggregation procedure.

(*Many*) F is *manipulable* at $(\mathcal{K}_1, \dots, \mathcal{K}_n)$ by individual i on p iff \mathcal{K}_i disagrees with $F(\mathcal{K}_1, \dots, \mathcal{K}_i, \dots, \mathcal{K}_n)$ on p , but \mathcal{K}_i agrees with $F(\mathcal{K}_1, \dots, \mathcal{K}_i^*, \dots, \mathcal{K}_n)$ on p for some alternative knowledge-base \mathcal{K}_i^* .

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- Hence both *PBP* and Fusion are manipulable on $\{C\}$. This seems bad.
- A possible defense (for both *PBP* and Fusion) is to say that the central cases for the application of these procedure are cases in which we need collective outcome + collective reasons.
- One could claim: in such cases strategic voting will not be attractive.

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- What is the supporter of Fusion to say in these cases? Does Fusion apply only when the need for collective reasons does not depend on any particular outcome on C ?
- If so, how are we aggregate judgments in cases where the need for collective reasons *does* depend on a specific outcome?
- If not, in what sense is manipulability on (sets containing) the conclusion not a problem for the Fusion approach?