

# Comments on *Judgment Aggregation without Paradox*

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May 25-29 2005 - FEW 2005 - UTA

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- Gabriella's arguments against PBP:
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  - **Connection**: applying majority-voting on isolated premises blinds the procedure to the logical connections between them. [Discussed in section 2]

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  - **Connection**: applying majority-voting on isolated premises blinds the procedure to the logical connections between them. [Discussed in section 2]
- One more common drawback of PBP & Fusion  
**Manipulability**: aggregated outcome can be manipulated by strategic voting. [Discussed in section 3]

The San Francisco Art School is looking for models (of any sex) for their painting classes. Jerry shows up for the auditions. A seven-member committee votes on the following:

- A: Jerry is attractive
- B: Jerry is poor in social skills
- C: Jerry is accepted
- **Integrity Constraint:**  $((A \& B) \vee (\sim A \& \sim B)) \equiv C$

	$A$	$B$	$C$	$A \& B$	$\sim A \& \sim B$	$A \equiv B$
Voter 1	1	1	1	1	0	1
Voter 2	1	1	1	1	0	1
Voter 3	1	1	1	1	0	1
Voter 4	1	0	0	0	0	0
Voter 5	1	0	0	0	0	0
Voter 6	0	1	0	0	0	0
Voter 7	0	1	0	0	0	0
Majority	1	1	0	0	0	0

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Voter 7	0	1	0	0	0	0
<b>Majority</b>	<b>1</b>	<b>1</b>	<b>0</b>	<b>0</b>	<b>0</b>	<b>0</b>

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- Example of *instability* of PBP: if we identify  $A$  and  $B$  as premises, PBP's outcome on  $C$  is 1.



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- Instance of doctrinal paradox.
- Example of *instability* of PBP: if we identify  $A$  and  $B$  as premises, PBP's outcome on  $C$  is 1.
- If we collect majority on the complex propositions  $A \& B$  and  $\sim A \& \sim B$  as premises, PBP's outcome on  $C$  is 0.

# Applying the Fusion Procedure

- $Mod(\mathcal{K}_1) = Mod(\mathcal{K}_2) = Mod(\mathcal{K}_3) = \{(1, 1, 1)\}$
- $Mod(\mathcal{K}_4) = Mod(\mathcal{K}_5) = \{(1, 0, 0)\}$
- $Mod(\mathcal{K}_6) = Mod(\mathcal{K}_7) = \{(0, 1, 0)\}$
- $Mod(IC) = \{(1, 1, 1), (0, 0, 1), (1, 0, 0), (0, 1, 0)\}$

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A	B	C	$\mathcal{K}_1$	$\mathcal{K}_2$	$\mathcal{K}_3$	$\mathcal{K}_4$	$\mathcal{K}_5$	$\mathcal{K}_6$	$\mathcal{K}_7$	$d_{\Sigma}(I, E)$
1	1	1	0	0	0	2	2	2	2	<b>8</b>
1	0	0	2	2	2	0	0	2	2	10
0	1	0	2	2	2	2	2	0	0	10
0	0	1	2	2	2	2	2	2	2	14

# A Language Variation

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- Let us now move to a language that does not contain  $A$ , and  $B$  but where there are atomic sentences  $D \approx (A \& B)$  and  $E \approx (\sim A \& \sim B)$ .
- Same vote as before, but now the integrity constraints are expressed in terms of  $D$ ,  $E$ , and  $C$ .
- **IC:**  $(D \vee E) \equiv C$  and  $D \rightarrow \sim E$ .
  - Note the extra-IC: it takes care of the connection between  $(A \& B)$  and  $(\sim A \& \sim B)$

# Reversal!

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$D$	$E$	$C$	$\mathcal{K}_1$	$\mathcal{K}_2$	$\mathcal{K}_3$	$\mathcal{K}_4$	$\mathcal{K}_5$	$\mathcal{K}_6$	$\mathcal{K}_7$	$d_{\Sigma}(I, E)$
1	0	1	0	0	0	2	2	2	2	8
0	1	1	2	2	2	2	2	2	2	14
0	0	0	2	2	2	0	0	0	0	<b>6</b>

# Instability and Linguistic Variance Compared

- By moving to a language in which the complex propositions  $(A \& B)$  and  $(\sim A \& \sim B)$  can be expressed as single atoms we reverse the outcome on the conclusion.



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  - (Proper) instability of PBP: different ways of applying the same procedure yield different outcomes—holding fixed the identification of the atoms.
  - In the current case, we need to carry out the fusion procedure exactly in the same way, but look at two different languages. (*linguistic instability*).
- The common feature between linguistic instability and instability proper is that in both cases the aggregation procedure sees a difference where there should be none.

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- Let  $\mathcal{L}_2$  have three sentences  $X, Y, Z$ , such that:  $Z : \approx C$ ,  
 $Y : \approx B \equiv C$ ,  $X : \approx A \equiv B$ .

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- Let  $\mathcal{L}_2$  have three sentences  $X, Y, Z$ , such that:  $Z : \approx C$ ,  $Y : \approx B \equiv C$ ,  $X : \approx A \equiv B$ .
- E.g.: A vote on  $(A \& B) \equiv C$  becomes a vote on  $[(X \equiv (Z \equiv Y)) \& (Z \equiv Y)] \equiv Z$ .

My conclusion in this section is that, to defend her approach, and at the same time be able to use *instability* as an objection to PBP, Gabriella owes both of the following:

- 1 An explanation of why *linguistic* instability is harmless for the Fusion procedure.



My conclusion in this section is that, to defend her approach, and at the same time be able to use *instability* as an objection to PBP, Gabriella owes both of the following:

- 1 An explanation of why *linguistic* instability is harmless for the Fusion procedure.
- 2 An account of why the explanation in (1) cannot be used by a supporter of PBP against the instability criticism.

- Distinction: (i) collective judgments in which the outcome must be supported by collectively accepted reasons vs. (ii) judgments in which it is enough to reach some collective conclusion.

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  - (ii) is better dealt by the Conclusion Based Procedure (CPB).

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- Standard view:
  - (i) is best dealt by procedures which, like PBP and fusion, give a complete outcome
  - (ii) is better dealt by the Conclusion Based Procedure (CPB).
- In the art-school case, we can imagine contexts in which the reasons for acceptance of  $C$  are important (e.g. if Jerry is accepted because attractive and unapproachable ( $A \& B$ ) then he gets a worse contract than he would get if he were accepted because unattractive and easygoing ( $\sim A \& \sim B$ ).

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  - In the art-school example, we had a majority for  $A$  and a majority for  $B$  but a minority for  $A \equiv B$ .
- But here, the fusion procedure recommends the same outcome as PBP, even though only a minority supports  $(A \& B)$  (and thus  $(A \equiv B)$ ).
- **Question:** in what sense does the Fusion procedure do better when it comes to keeping track of logical connections between premises?



# Independence on a Set of Propositions

- F. Dietrich and C. List, *Strategy-proof judgment aggregation*, forthcoming.

(*Ind<sub>Y</sub>*) An aggregation function  $F$  is independent (on a set of propositions  $Y$ ) iff for every proposition  $p \in Y$ , and  $n \in \omega$  there is a function  $\phi : \{0, 1\}^n \rightarrow \{0, 1\}$  such that  $p \in F(\mathcal{K}_1, \dots, \mathcal{K}_m)$  iff  $\phi(v_1(p), \dots, v_n(p)) = 1$  where  $v_i(p)$  is  $\mathcal{K}_i$ 's vote on  $p$ .

- For each member  $p$  of  $Y$  the aggregated outcome on  $p$  is a function of the pattern of individual judgments on  $p$ .

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- For each member  $p$  of  $Y$  the aggregated outcome on  $p$  is a function of the pattern of individual judgments on  $p$ .
- In our example, whenever  $Y \supset \{C\}$ , both *PBP* and *Fusion* are not Independent on  $Y$ .

	A	B	C		A	B	C
Voter 1	1	1	<b>1</b>		1	1	<b>1</b>
Voter 2	1	1	<b>1</b>		1	1	<b>1</b>
Voter 3	1	1	<b>1</b>		1	1	<b>1</b>
Voter 4	1	0	<b>0</b>		1	0	<b>0</b>
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Voter 7	0	1	<b>0</b>		1	0	<b>0</b>
<b>PBP + Fusion</b>	<b>1</b>	<b>1</b>	<b>1</b>		<b>1</b>	<b>0</b>	<b>0</b>

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- Strategic voting on  $p$  is voting untruthfully on other propositions so as to change the collective outcome on  $p$  into the desired outcome.
- More formally, let  $F$  be an aggregation procedure.

(*Many*)  $F$  is *manipulable* at  $(\mathcal{K}_1, \dots, \mathcal{K}_n)$  by individual  $i$  on  $p$  iff  $\mathcal{K}_i$  disagrees with  $F(\mathcal{K}_1, \dots, \mathcal{K}_i, \dots, \mathcal{K}_n)$  on  $p$ , but  $\mathcal{K}_i$  agrees with  $F(\mathcal{K}_1, \dots, \mathcal{K}_i^*, \dots, \mathcal{K}_n)$  on  $p$  for some alternative knowledge-base  $\mathcal{K}_i^*$ .

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- Hence both *PBP* and Fusion are manipulable on  $\{C\}$ . This seems bad.
- A possible defense (for both *PBP* and Fusion) is to say that the central cases for the application of these procedure are cases in which we need collective outcome + collective reasons.
- One could claim: in such cases strategic voting will not be attractive.

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- What is the supporter of Fusion to say in these cases? Does Fusion apply only when the need for collective reasons does not depend on any particular outcome on  $C$ ?
- If so, how are we aggregate judgments in cases where the need for collective reasons *does* depend on a specific outcome?
- If not, in what sense is manipulability on (sets containing) the conclusion not a problem for the Fusion approach?